TOA

LECTURE 6

Muhammad Hafeez Javed Foundation University Rawalpindi rmhjaved@gmail.com www.rmhjaved.com

Recap lecture 5

Different notations of transition diagrams, languages of strings of even length, Odd length, starting with b, ending in a (with different FAs), beginning with a, not beginning with b, beginning with and ending in same letters

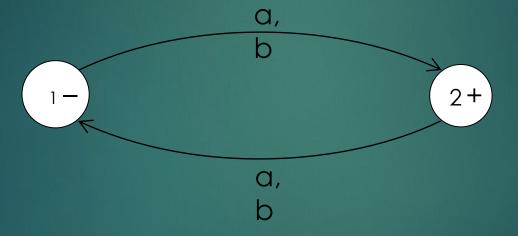
TASK

Build an FA for the language L of strings, defined over $\Sigma = \{a, b\}$, **Of** odd length.

Solution: The language L may be expressed by RE (a+b)((a+b)(a+b))* or ((a+b)(a+b))*(a+b)

This language may be accepted by the following FA

Solution continued ...



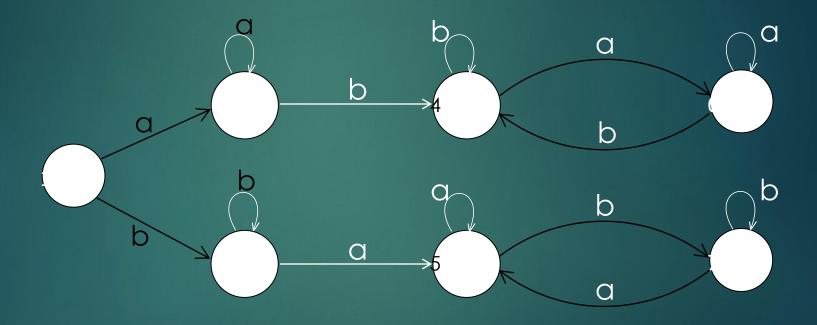
▶ Build an FA accepting the Language L of Strings, defined over $\Sigma = \{a, b\}$, beginning with and ending in same letters.

Solution:The language L may be expressed by the following regular expression

$$(a+b)+a(a+b)^*a+b(a+b)^*b$$

This language L may be accepted by the following FA

Solution continued ...

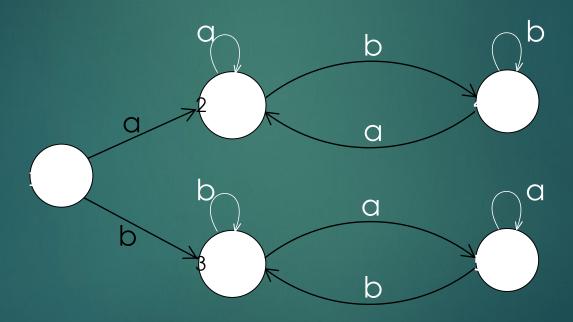


Consider the Language L of Strings , defined over $\Sigma = \{a, b\}$, beginning with and ending in different letters.

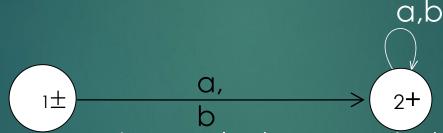
The language L may be expressed by the following regular expression

$$a (a + b)^* b + b (a + b)^* a$$

This language may be accepted by the following FA



Consider the Language L, defined over Σ = {a, b} of all strings including Λ, The language L may be accepted by the following FA



The language L may also be accepted by the following FA



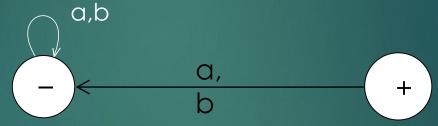
The language L may be expressed by the following regular expression

$$(a + b)^*$$

Consider the Language L, defined over Σ = {a, b} of all non empty strings. The language L may be accepted by the following FA

The above language may be expressed by the following regular expression (a + b)⁺

► Consider the following FA, defined over $\Sigma = \{a, b\}$

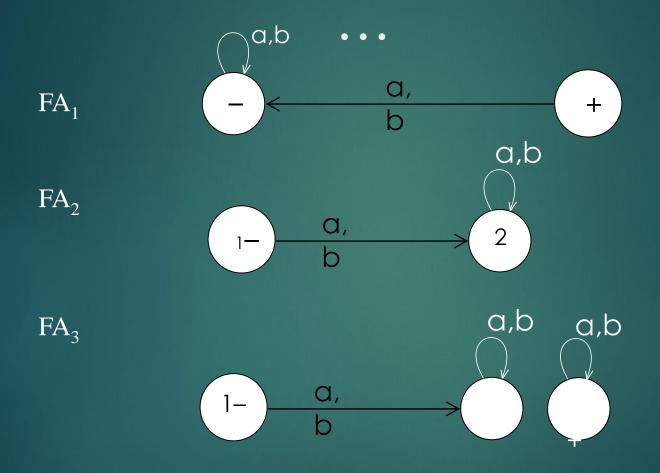


▶ It is to be noted that the above FA does not accept any string. Even it does not accept the null string. As there is no path starting from initial state and ending in final state.

Equivalent FAs

It is to be noted that two FAs are said to be equivalent, if they accept the same language, as shown in the following FAs.

Equivalent FAs Continued



Note (Equivalent FAs)

► FA₁ has already been discussed, while in FA₂, there is no final state and in FA₃, there is a final state but FA₃ is disconnected as the states 2 and 3 are disconnected.

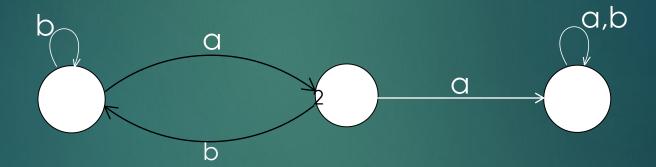
It may also be noted that the language of strings accepted by FA_1 , FA_2 and FA_3 is denoted by the empty set *i.e.*

 $\{\}$ OR \emptyset

Consider the Language L of strings, defined over $\Sigma = \{a, b\}$, containing double a.

The language L may be expressed by the following regular expression

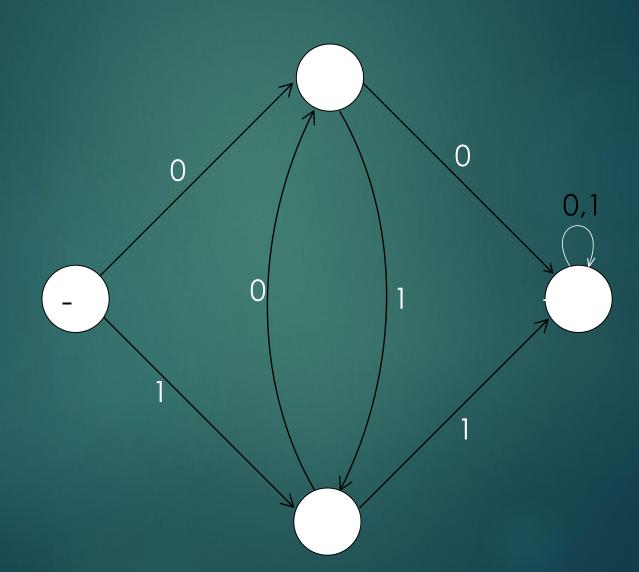
(a+b)* (aa) (a+b)*. This language may be accepted by the following FA



Consider the language L of strings, defined over

 Σ ={0, 1}, **having double 0's or double 1's**, The language L may be expressed by the regular expression $(0+1)^*(00+11)(0+1)^*$

This language may be accepted by the following FA

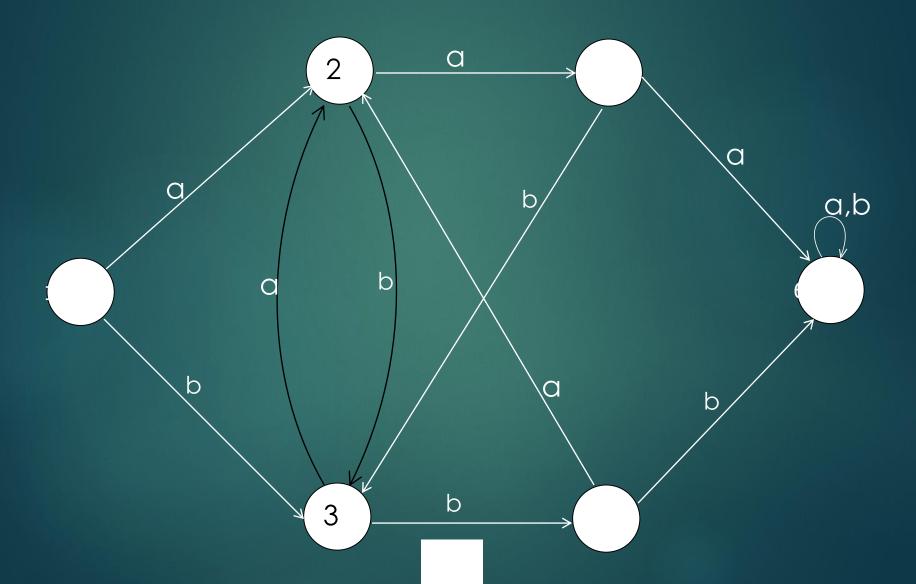


Consider the language L of strings, defined over

 $\Sigma = \{A, b\}$, having triple a's or triple b's. The language L may be expressed by RE

 $(a+b)^*$ (aaa + bbb) $(a+b)^*$

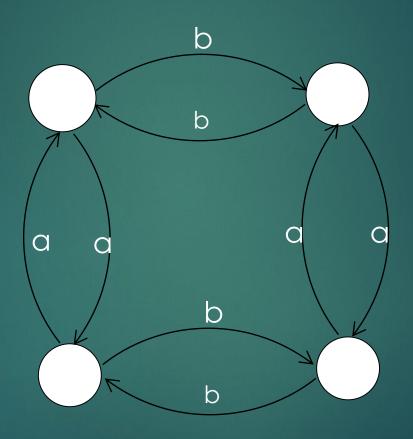
This language may be accepted by the following FA



Consider the EVEN-EVEN language, defined over Σ={a, b}. As discussed earlier that EVEN-EVEN language can be expressed by the regular expression

 $(aa+bb+(ab+ba)(aa+bb)^*(ab+ba))^*$

EVEN-EVEN language may be accepted by the following FA



Summing Up

Language of strings beginning with and ending in different letters, Accepting all strings, accepting non-empty strings, accepting no string, containing double a's, having double 0's or double 1's, containing triple a's or triple b's, EVEN-EVEN